Free Rota-Baxter family algebras and free (tri)dendriform family algebras

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joint work with Xing Gao and Dominique Manchon

Algebraic, analytic and geometric structures emerging from quantum field theory, ${\sf Chengdu}$

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Outline

- Free Rota-Baxter family algebras (RBFA)
 - Typed angularly decorated planar rooted trees
 - A multiplication on $\mathbf{k}\mathcal{T}(X,\Omega)$
 - Construction of free RBFA
- Embedding free DFAs (resp.TFAs) into free RBFAs
 - Embedding free DFAs into free RBFAs
 - Embedding free TFAs into free RBFAs
 - Universal enveloping algebras of (tri)dendriform family algebras

- In 1998, Loday and Ronco proved the free dendriform algebra on one generator can be described as an algebra over the set of planar binary trees.
- In 2004, Loday and Ronco showed the free tridendriform algebra on one generator can be described as an algebra over the set of planar trees.
- J. -L. Loday and M. O. Ronco, Hopf algebra of the planar binary trees, *Adv. Math.* **39** (1998), 293-309.
- J. -L. Loday and M. O. Ronco, Trialgebras and families of polytopes, in "Homotopy theoty: relations with algebraic geometry, group cohomology, and algebraic K-theory", *Contemp. Math.* **346** (2004), 369-398.

- In 2007, K. Ebrahimi-Fard, J. Gracia-Bondia and F. Patras studied about algebraic aspects of renormalization in Quantum field theory. The first example about Rota-Baxter family algebras of -1 appeared in this paper and Guo suggested them to use this name.
- In 2008, K. Ebrahimi-Fard and L. Guo, they used rooted trees and forests to give explicit constructions of free noncommutative Rota-Baxter algebras on modules and sets.
- K. Ebrahimi-Fard, J. Gracia-Bondia and F. Patras, A Lie theoretic approach to renormalization, *Comm. Math. Phys.* **276** (2007), 519-549.
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- In 2009, L. Guo named Rota-Baxter family algebras of weight λ .
- In 2012, E. Panzer studied algebraic aspects of renormalization in Quantum Field Theory. They proved the Taylor expansion operators fulfil for Rota-Baxter family algebras of weight -1.
- L. Guo, Operated monoids, Motzkin paths and rooted trees, *J. Algebraic Combin.* **29** (2009), 35-62.
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- In 2018, Bruned, Haier and Zambotti gave a systematic description of a canonical renormalisation procedure of stochastic PDEs, in which their construction is based on bialgebras of typed decorated forests in cointeraction.
- In 2018, L. Foissy studied multiple prelie algebras and related operads. He proved the free *T*-multiple prelie algebra generated by a set *D*.
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Questions

- How to construct the free Rota-Baxter family algebras by using typed angularly decorated planar rooted trees?
- Whow to construct the free (tri)dendriform family algebra by using typed decorated planar (binary) rooted trees?
- What's the relationship between the free Rota-Baxter family algebra and the free (tri)dendriform family algebra?

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Definition and Example

Definition (Ebrahimi-Fard et al. 2007, Guo 2009)

Let Ω be a semigroup and $\lambda \in \mathbf{k}$ be given. A **Rota-Baxter family** of weight λ on an algebra R is a collection of linear operators $(P_{\omega})_{\omega \in \Omega}$ on R such that

$$P_{\alpha}(a)P_{\beta}(b) = P_{\alpha\beta}(P_{\alpha}(a)b + aP_{\beta}(b) + \lambda ab),$$

where $a, b \in R$ and $\alpha, \beta \in \Omega$. Then the pair $(R, (P_{\omega})_{\omega \in \Omega})$ is called a **Rota-Baxter family algebra** of weight λ .

Example 1

The algebra of Laurent series $R=\mathbf{k}\left[z^{-1},z\right]$ is a Rota-Baxter family algebra of weight -1 , with $\Omega=(\mathbb{Z},+)$, where the operator P_{ω} is the projection onto the subspace $R_{<\omega}$ generated by $\left\{z^k,k<\omega\right\}$ parallel to the supplementary subspace $R_{\geq\omega}$ generated by $\left\{z^k,k\geq\omega\right\}$.

Example 2

Let $\Omega=(\mathbb{R},+)$ be a semigroup, and R is the \mathbb{R} -algebra consisting of all continous functions from \mathbb{R} to \mathbb{R} , the multiplication is defined by

$$(fg)(x) := f(x)g(x), \text{ for } f,g \in R.$$

And for any $\alpha \in \Omega$, define a family linear operator $P_{\alpha}: R \to R$ as follows

$$P_{\alpha}(f)(x) = e^{-\alpha A(x)} \int_{0}^{x} e^{\alpha A(t)} f(t) dt$$
, for $\alpha \in \Omega$,

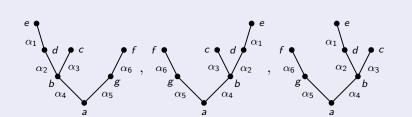
where A is a fixed nonzero element of R. Then $(P_{\alpha})_{\alpha \in \Omega}$ is a Rota-Baxter family of weight 0.

Definition (Bruned-Hairer-Zambotti, Invent. math.)

Let X and Ω be two sets. An X-decorated Ω -typed (abbreviated typed decorated) rooted tree is a triple T = (T, dec, type), where

- T is a rooted tree.
- $ext{ dec}: V(T) \to X ext{ is a map.}$
- **3** type : $E(T) \rightarrow \Omega$ is a map.

Example 3

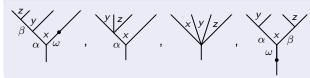


Definition

Let X and Ω be two sets. An X-angularly decorated Ω -typed (abbreviated typed angularly decorated) planar rooted tree is a triple T = (T, dec, type), where

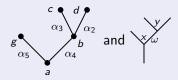
- ① T is a planar rooted tree.
- $\mathbf{2} \ \mathrm{dec} : A(T) \to X \ \mathrm{is \ a \ map}.$
- **3** type : $IE(T) \rightarrow \Omega$ is a map.

Example 4



Remark

The graphical representation of (planar) rooted trees



in Example 4 and Example 3 is slightly different.

- Here the root and the leaves are now edges rather than vertices.
- The set E(T) must be replaced by the set IE(T) of internal edges.

If a semigroup Ω has no identity element, we consider the monoid $\Omega^1:=\Omega\sqcup\{1\}$ obtained from Ω by adjoining an identity:

$$1\omega := \omega 1 := \omega, \text{ for } \omega \in \Omega \text{ and } 11 := 1.$$

For $n \geq 0$, let $\mathcal{T}_n(X, \Omega)$ denote the set of X-angularly decorated Ω^1 -typed planar rooted trees with n+1 leaves such that leaves are decorated by the identity 1 in Ω^1 and internal edges are decorated by elements of Ω . Note that the root is not decorated. Denote by

$$\mathcal{T}(X,\,\Omega):=\bigsqcup_{n\geq 0}\mathcal{T}_n(X,\,\Omega)\ \ \text{and}\ \ \mathbf{k}\mathcal{T}(X,\,\Omega):=\bigoplus_{n\geq 0}\mathbf{k}\mathcal{T}_n(X,\,\Omega).$$

Graphically, an element $T \in \mathcal{T}(X, \Omega)$ is of the form:

$$T = T_1 \underbrace{\bigcap_{\alpha_1 \times 1}^{T_2} T_n}_{\alpha_1 \times 1} T_{n+1}, \text{ with } n \ge 0,$$

i.e., $\alpha_j\in\Omega$ if $T_j\neq |;\ \alpha_j=1$ if $T_j=|.$ For each $\omega\in\Omega,$ define a linear operator

$$B_{\omega}^{+}: \mathbf{k}\mathcal{T}(X, \Omega) \to \mathbf{k}\mathcal{T}(X, \Omega),$$

by adding a new root and a new internal edge decorated by ω connecting the new root and the root of T. For example,

$$B_{\omega}^{+}\Big(\Big|\Big) = \left|_{\omega}, \quad B_{\omega}^{+}\Big(\bigvee^{\times}\Big) = \bigcup^{\times}_{\omega}, \quad B_{\omega}^{+}\Big(\bigvee^{\times}_{\alpha}\bigvee^{\times}_{\beta}\Big) = \bigcup^{\times}_{\omega}\bigvee^{\times}_{\beta}.$$

The **depth** dep(T) of a rooted tree T is the maximal length of linear chains from the root to the leaves of the tree. For example,

$$dep(\left| \begin{array}{c} \\ \end{array} \right|) = dep\left(\begin{array}{c} \\ \\ \end{array} \right) = 1 \text{ and } dep\left(\begin{array}{c} \\ \\ \end{array} \right) = 2.$$

We add the "zero-vertex tree" | to the picture, and set dep(|)=0. Note that the operators B_{ω}^+ are not defined on |.

Remark

For any $T \in \mathcal{T}(X, \Omega) \sqcup \{|\}$, if $\operatorname{dep}(T) = 0$, then T = |. Define the number of branches $\operatorname{bra}(T)$ of T to be 0 in this case. Otherwise, $\operatorname{dep}(T) \geq 1$ and T is of the form

$$T = T_1 \underbrace{\bigcap_{\alpha_1 \times 1}^{T_2} T_n}_{\alpha_1 \times \alpha_{n+1}} T_{n+1} \text{ with } n \geq 0.$$

Here any branch $T_j \in \mathcal{T}(X, \Omega) \sqcup \{|\}, j = 1, \ldots, n+1$ is of depth at most one less than the depth of T, and equal to zero if and only if $T_j = |$. We define $\operatorname{bra}(T) := n+1$. For example,

$$\operatorname{bra}\left(\begin{array}{c} \omega \end{array}\right) = 1, \ \operatorname{bra}\left(\begin{array}{c} x \\ \end{array}\right) = 2 \ \text{and} \ \operatorname{bra}\left(\begin{array}{c} x \\ \end{array}\right) = 3.$$

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We define $T \diamond T'$ by induction on $\operatorname{dep}(T) + \operatorname{dep}(T') \geq 2$. For the initial step $\operatorname{dep}(T) + \operatorname{dep}(T') = 2$, we have $\operatorname{dep}(T) = \operatorname{dep}(T') = 1$ and T, T' are of the form

$$T = \underbrace{\begin{array}{c} x_1 \\ \\ \end{array}}_{x_m} \text{ and } T' = \underbrace{\begin{array}{c} y_1 \\ \\ \end{array}}_{y_n}, \text{ with } m, n \geq 0.$$

Define

$$T \diamond T' := \underbrace{x_1 \cdots x_m}_{x_1 \cdots x_m} \diamond \underbrace{y_1 \cdots y_n}_{y_1 \cdots y_n} := \underbrace{x_1 \cdots x_m}_{y_1 \cdots y_n}.$$

For the induction step $dep(T) + dep(T') \ge 3$, the trees T and T' are of the form

$$T = T_1 \underbrace{\bigcirc_{\alpha_1}^{T_2} T_m}_{\alpha_1} \underbrace{\bigcirc_{\alpha_2 \alpha_m \circ}^{T_2}}_{T_{m+1}} T_{m+1} \text{ and } T' = \underbrace{T'_1}_{\beta_1} \underbrace{\bigcirc_{\beta_1}^{T_2}}_{\beta_{n+1}} \underbrace{\neg_{n+1}^{T_2}}_{\beta_{n+1}} T'_{n+1}.$$

There are four cases to consider.

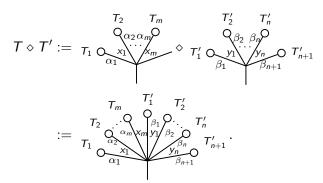
Case 1: $T_{m+1} = | = T'_1$. Define

Case 2: $T_{m+1} \neq | = T'_1$. Define

$$T \diamond T' := T_{1} \underbrace{ \begin{array}{c} T_{2} & T_{m} & T'_{2} & T'_{n} \\ \alpha_{2} \alpha_{m} & \alpha_{m+1} \end{array}}_{X_{1} & X_{m+1}} T_{m+1} \diamond \underbrace{ \begin{array}{c} T_{2} & T'_{n} \\ \beta_{2} & \beta_{n} \end{array}}_{Y_{n+1}} T'_{n+1}$$

$$:= \underbrace{ \begin{array}{c} T_{m} & T_{m+1} & T'_{2} \\ T_{2} & \vdots & \alpha_{m} & T'_{n} \end{array}}_{T_{1} & \beta_{n+1}} T'_{n+1} \cdot \underbrace{ \begin{array}{c} T_{m} & T_{m+1} \\ \gamma_{1} & \vdots \\ \gamma_{n} & \alpha_{m} & \gamma_{1} & \beta_{n} \end{array}}_{X_{1} & \beta_{n+1}} T'_{n+1} \cdot \underbrace{ \begin{array}{c} T_{m} & T_{m+1} \\ \gamma_{1} & \vdots \\ \gamma_{n} & \alpha_{1} & \vdots \end{array}}_{X_{1} & \beta_{n+1}} T'_{n+1} \cdot \underbrace{ \begin{array}{c} T_{m} & T_{m+1} \\ \gamma_{1} & \vdots \\ \gamma_{n} & \beta_{n+1} \end{array}}_{X_{1} & \vdots } T'_{n+1} \cdot \underbrace{ \begin{array}{c} T_{m} & T_{m+1} \\ \gamma_{1} & \vdots \\ \gamma_{n} & \vdots \\ \gamma$$

Case 3: $T_{m+1} = | \neq T'_1$. Define



Case 4: $T_{m+1} \neq | \neq T'_1$. Define

$$T \diamond T' := T_{1} \underbrace{ \begin{array}{c} T_{2} & T_{m} & T'_{2} & T'_{n} \\ \alpha_{2} \alpha_{m} & \beta_{2} & \beta_{n} \\ \end{array}}_{\alpha_{m+1}} T_{m+1} \diamond T'_{1} \underbrace{ \begin{array}{c} \beta_{2} & \beta_{n} \\ \beta_{1} & \beta_{n+1} \\ \end{array}}_{\beta_{n+1}} T'_{n+1}$$

$$:= \left(T_{1} \underbrace{ \begin{array}{c} T_{2} & T_{m} \\ \alpha_{2} \alpha_{m} \\ \end{array}}_{\alpha_{1}} \diamond \left(B_{\alpha_{m+1}}^{+}(T_{m+1}) \diamond B_{\beta_{1}}^{+}(T'_{1}) \right) \right) \diamond \underbrace{ \begin{array}{c} T'_{2} & T'_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n+1}} T'_{n+1}$$

$$:= \left(T_{1} \underbrace{ \begin{array}{c} T_{2} & T_{m} \\ \alpha_{2} \alpha_{m} \\ \end{array}}_{\alpha_{1}} \diamond B_{\alpha_{m+1}}^{+} \beta_{1} \left(B_{\alpha_{m+1}}^{+}(T_{m+1}) \diamond T'_{1} + T_{m+1} \diamond B_{\beta_{1}}^{+}(T'_{1}) + \lambda T_{m+1} \diamond T'_{1} \right) \right)$$

$$T'_{2} & T'_{n} \\ \underbrace{ \begin{array}{c} \beta_{2} & \beta_{n} \\ \\ \beta_{1} & \beta_{n} \\ \end{array}}_{\beta_{1}} \underbrace{ \begin{array}{c} T_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ \beta_{2} & \beta_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ T_{n} \\ \end{array}}_{\beta_{n}} \underbrace{ \begin{array}{c} T_{n} \\ T_{n}$$

Example 6

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Problems

- About the Rota-Baxter family algebras on typed-angularly decorated planar rooted trees constructed above: are they free?
- What's the relationship between typed-angularly decorated planar rooted trees and bracketed words?

Free FRBA on bracketed words

Denote by S the following vector subspace of $\mathbf{k}\mathfrak{M}(\Omega,X)$:

$$S := \left\{ \left\lfloor x \right\rfloor_{\alpha} \left\lfloor y \right\rfloor_{\beta} - \left\lfloor \left\lfloor x \right\rfloor_{\alpha} y \right\rfloor_{\alpha\beta} - \left\lfloor x \left\lfloor y \right\rfloor_{\beta} \right\rfloor_{\alpha\beta} - \lambda \left\lfloor xy \right\rfloor_{\alpha\beta} \right\},\,$$

where $\alpha, \beta \in \Omega, x, y \in \mathfrak{M}(\Omega, X)$.

Theorem (Z.-Gao, Pacific J. Math.)

Let X be a set and let Ω be a semigroup.

- S is a Gröbner-Shirshov basis in $\mathbf{k}\mathfrak{M}(\Omega,X)$ with respect to a monomial order.
- ② The set \mathfrak{X}_{∞} is a **k**-basis of the free Rota-Baxter family algebra $(\mathbf{k}X_{\infty}, \diamond_w, (\lfloor \rfloor \omega)_{\omega \in \Omega}) = \mathbf{k}\mathfrak{M}(\Omega, X)/\operatorname{Id}(S)$ of weight λ , where $\operatorname{Id}(S)$ is the operated ideal generated by S in $\operatorname{kM}(\Omega, X)$.
- Y. Y. Zhang and X. Gao, Free Rota-Baxter family algebras and (tri)dendriform family algebras, *Pacific J. Math.* **301** (2019), 741-766.

The relationship

Built up a one-one correspondence between $\mathbf{k}\mathcal{T}(X,\Omega)$ and $\mathbf{k}\mathfrak{X}_{\infty}$.

| Typed angularly decorated planar rooted trees $\mathcal{T}(X,\Omega)$ | \mathfrak{X}_{∞} |
|---|------------------------------|
| | 1 |
| ω | $\lfloor 1 floor_\omega$ |
| × | Х |
| w X | $\lfloor x \rfloor_{\omega}$ |
| × | xy |
| x y w | $[xy]_{\omega}$ |

Construction of the isomorphism map

Define a linear map

$$\phi: \mathbf{k}\mathcal{T}(X, \Omega) \to \mathbf{k}\mathfrak{X}_{\infty}, \quad T \mapsto \phi(T)$$

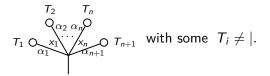
by induction on $dep(T) \ge 1$. If dep(T) = 1, then $T = x_1$ with n > 0 and define

$$\phi(T) := \phi(\underbrace{x_1 \cdots x_n}_{x_n}) := x_1 \cdots x_n,$$

with $x_1 \cdots x_n := 1$ in the case n = 0.

Construction of the isomorphism map

For the induction step $dep(T) \ge 2$, T is the form



Then we define $\phi(T)$ by the induction on $\operatorname{bra}(T) \geq 1$. For the initial step $\operatorname{bra}(T) = 1$, we have $T = \bigcap_{\alpha_1}^{T_1}$ and define

$$\phi\left(\bigcap_{\alpha_1}^{I_1}\right) := \phi(B_{\alpha_1}^+(T_1)) := \lfloor \phi(T_1) \rfloor_{\alpha_1}.$$

Construction of the isomorphism map

For the induction step $bra(T) \ge 2$, there are two cases to consider.

Case 1: $T_1 = |$. Define

$$\phi(T) := x_1 \phi\left(\begin{array}{ccc} T_3 & T_n \\ \alpha_3 & \alpha_n \\ T_2 & \alpha_2 & x_2 \\ & & x_n \\ & & x_{n+1} \end{array}, T_{n+1} \right).$$

Case 2: $T_1 \neq |$. Define

$$\phi(T) := \lfloor \phi(T_1) \rfloor_{\alpha_1} x_1 \phi \left(\begin{array}{ccc} T_3 & T_n \\ Q_{\alpha_3 & \alpha_n} Q \\ T_2 & X_2 & X_n \\ Q_{\alpha_2} & X_{n+1} \end{array} \right).$$

Construction of the isomorphism map

Conversely, define a linear map

$$\psi: \mathbf{k} X_{\infty} \to \mathbf{k} \mathcal{T}(X, \Omega), \quad w \mapsto \psi(w)$$

by induction on $dep(w) \ge 1$.

- If dep(w) = 1, then $w = x_1 \cdots x_n \in M(X)$ with $n \ge 0$.
- If dep(w) > 1, we apply induction on $bre(w) \ge 1$.

Write $w=w_1\cdots w_n$ with $\operatorname{bre}(w)=n\geq 1$. If $\operatorname{bre}(w)=1$, then $w=\lfloor \overline{w}\rfloor_{\alpha}$ for $\overline{w}\in\mathfrak{X}_{\infty}$ and $\alpha\in\Omega$ by

 $dep(w) \ge 1$, and define

$$\psi(\mathbf{w}) = \psi\left(\lfloor \overline{\mathbf{w}} \rfloor_{\alpha}\right) := B_{\alpha}^{+}(\psi(\overline{\mathbf{w}})).$$

If $bre(w) \ge 2$, then define

$$\psi(\mathbf{w}) := \psi(\mathbf{w}_1) \diamond \psi(\mathbf{w}_2 \cdots \mathbf{w}_n).$$

Construction of the isomorphism map

Proposition (Z.-Gao-Manchon, Algebr. Represent. Theory) We have $\psi \circ \phi = id$ and $\phi \circ \psi = id$.

Lemma (Z.-Gao-Manchon, Algebr. Represent. Theory)

For T and T' in $\mathcal{T}(X,\Omega)$, we have

$$\phi(T \diamond T') = \phi(T) \diamond_{\mathsf{w}} \phi(T').$$



Y. Y. Zhang, X. Gao and D. Manchon, Free Rota-Baxter family algebras and free (tri)dendriform family algebras, Algebr. Represent. Theory, 301 (2023), 741-766.

Main results

Let j_X be the embedding given by

$$j_X: X \to \mathbf{k}\mathcal{T}(X,\Omega), x \mapsto X$$

Theorem (Z.-Gao-Manchon, Algebr. Represent. Theory)

Let X be a set and let Ω be a semigroup. The triple $(\mathbf{k}\mathcal{T}(X,\Omega),\diamond,(B^+_\omega)_{\omega\in\Omega})$, together with the j_X , is the free Rota-Baxter family algebra of weight λ on X.

Corollary (Z.-Gao-Manchon, Algebr. Represent. Theory)

Let X be a set and let Ω be a trivial semigroup with one element. Then the triple $(\mathbf{k}\mathcal{T}(X), \diamond, B^+)$, together with the j_X , is the free Rota-Baxter algebra of weight λ on X.

- 1 Free Rota-Baxter family algebras (RBFA)
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Problems

- How to construct free (tri)dendriform family algebras via typed decorated planar rooted trees?
- What's the relationship between free Rota-Baxter family algebras and free (tri)dendriform family algebras?

Definition

Let Ω be a semigroup. A **dendriform family algebra** is a **k**-module D with a family of binary operations $(\prec_{\omega}, \succ_{\omega})_{\omega \in \Omega}$ such that for $x, y, z \in D$ and $\alpha, \beta \in \Omega$,

$$(x \prec_{\alpha} y) \prec_{\beta} z = x \prec_{\alpha\beta} (y \prec_{\beta} z + y \succ_{\alpha} z),$$

$$(x \succ_{\alpha} y) \prec_{\beta} z = x \succ_{\alpha} (y \prec_{\beta} z),$$

$$(x \prec_{\beta} y + x \succ_{\alpha} y) \succ_{\alpha\beta} z = x \succ_{\alpha} (y \succ_{\beta} z).$$

For $n \ge 1$, let

$$Y_n := Y_{n,X,\Omega} := \mathcal{T}_n(X,\Omega) \cap \{\text{planar binary trees}\}.$$

Example 7
$$Y_{1} = \left\{ \begin{array}{c} x \\ \\ \end{array} \middle| x \in X \right\},$$

$$Y_{2} = \left\{ \begin{array}{c} x \\ \\ \end{array} \middle| x, y \in X, \alpha \in \Omega \right\},$$

$$Y_{3} = \left\{ \begin{array}{c} x \\ \\ \end{array} \middle| x, y \in X, \alpha \in \Omega \right\},$$

$$Y_{3} = \left\{ \begin{array}{c} x \\ \\ \end{array} \middle| x, y \in X, \alpha \in \Omega \right\},$$

The grafting $\vee_{x,(\alpha,\beta)}$ over x and (α,β) is defined to be $T=T^I\vee_{x,(\alpha,\beta)}T^r$ for some $x\in X$ and $\alpha,\beta\in\Omega^1$.

Example 8

$$\begin{array}{c}
\stackrel{y}{\underset{\alpha}{\bigvee}} = & \stackrel{y}{\underset{\alpha}{\bigvee}} \vee_{x,(\alpha,1)} |, & \stackrel{y}{\underset{\alpha}{\bigvee}} = | \vee_{x,(1,\alpha)} & \stackrel{y}{\underset{\alpha}{\bigvee}}, \\
\stackrel{y}{\underset{\alpha}{\bigvee}} \stackrel{z}{\underset{\beta}{\bigvee}} = & \stackrel{y}{\underset{\alpha}{\bigvee}} \vee_{x,(\alpha,\beta)} & \stackrel{z}{\underset{\alpha}{\bigvee}}.
\end{array}$$

Denote by

$$\mathrm{DD}(X,\Omega) := \bigoplus_{n \geq 1} \mathbf{k} Y_n.$$

Definition

Let X be a set and let Ω be a semigroup. Define binary operations

$$\prec_{\omega}, \succ_{\omega}: \left(\mathrm{DD}(X, \Omega) \otimes \mathrm{DD}(X, \Omega) \right) \oplus \left(\mathbf{k} | \otimes \mathrm{DD}(X, \Omega) \right)$$
$$\oplus \left(\mathrm{DD}(X, \Omega) \otimes \mathbf{k} | \right) \to \mathrm{DD}(X, \Omega), \text{ for } \omega \in \Omega$$

recursively on dep(T) + dep(U) by

- ② For $T = T^I \vee_{x,(\alpha_1,\alpha_2)} T^r$ and $U = U^I \vee_{y,(\beta_1,\beta_2)} U^r$, define

$$T \prec_{\omega} U := T^{I} \vee_{x, (\alpha_{1}, \alpha_{2}\omega)} (T^{r} \prec_{\omega} U + T^{r} \succ_{\alpha_{2}} U),$$

$$T \succ_{\omega} U := (T \prec_{\beta_{1}} U^{I} + T \succ_{\omega} U^{I}) \vee_{y, (\omega\beta_{1}, \beta_{2})} U^{r}.$$

Remark

Note that $|\prec_{\omega}|$ and $|\succ_{\omega}|$ are not defined for $\omega\in\Omega^1$. Here we apply the convention that

$$|\succ_1 T := T \prec_1 | := T$$
 and $|\prec_1 T := T \succ_1 | := 0$.

Example 9

Let
$$T = \bigvee_{x} \text{ and } U = \bigvee_{y} \text{ with } x, y \in X. \text{ For } \omega \in \Omega,$$

$$\bigvee_{x} \swarrow_{\omega} \bigvee_{y} = (|\vee_{x,(1,1)}|) \prec_{\omega} \bigvee_{y} + |\succ_{1} \bigvee_{y})$$

$$= |\vee_{x,(1,\omega)} \bigvee_{y} + |\succ_{1} \bigvee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} + |\succ_{1} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} + |\sim_{1} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} + |\sim_{1} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} + |\sim_{1} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y,(1,\omega)} |\sim_{1} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} + |\sim_{1} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} \vee_{y} + |\sim_{1} \vee_{y} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} \vee_{y} + |\sim_{1} \vee_{y} \vee_{y} \rangle$$

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$$= |\vee_{x,(1,\omega)} \vee_{y} \vee_{y} \vee_{y} \vee_{y} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} \vee_{y} \vee_{y} \vee_{y} \rangle$$

$$= |\vee_{x,(1,\omega)} \vee_{y} \vee_{y}$$

Let $j: X \to \mathrm{DD}(X,\Omega)$ be the natural embedding map defined by $j(x) = \bigvee_{x \in X} for \ x \in X.$

Theorem (Z.-Gao-Manchon, Algebr. Represent. Theory)

Let X be a set and let Ω be a semigroup. Then $(\mathrm{DD}(X,\Omega),(\prec_{\omega},\succ_{\omega})_{\omega\in\Omega})$, together with the map j, is the free dendriform family algebra on X.

Lemma (Z.-Gao, Pacific J. Math.)

Let X be a set and let Ω be a semigroup. The Rota-Baxter family algebra $(\mathbf{k}\mathcal{T}(X,\Omega),\diamond,(B_{\omega}^+)_{\omega\in\Omega})$ of weight ${\color{blue}0}$ induces a dendriform family algebra $(\mathbf{k}\mathcal{T}(X,\Omega),(\prec_{\omega}',\succ_{\omega}')_{\omega\in\Omega})$, where

$$T \prec_{\omega}' U := T \diamond B_{\omega}^+(U) \text{ and } T \succ_{\omega}' U := B_{\omega}^+(T) \diamond U,$$

with $T, U \in \mathcal{T}(X, \Omega)$.

Now, we give the relationship between the free dendriform family algebra and the free Rota-Baxter family algebra.

Theorem (Z.-Gao-Manchon, Algebr. Represent. Theory)

Let X be a set and let Ω be a semigroup. The free dendriform family algebra $(\mathrm{DD}(X,\Omega),(\prec_{\omega},\succ_{\omega})_{\omega\in\Omega})$ on X is a dendriform family subalgebra of the free Rota-Baxter family algebra $(\mathbf{k}\mathcal{T}(X,\Omega),\diamond,(B^+_{\omega})_{\omega\in\Omega})$ of weight 0.

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Definition

Let Ω be a semigroup. A **tridendriform family algebra** is a **k**-module T equipped with a family of binary operations $(\prec_{\omega}, \succ_{\omega})_{\omega \in \Omega}$ and a binary operation \cdot such that for $x, y, z \in T$ and $\alpha, \beta \in \Omega$,

$$(x \prec_{\alpha} y) \prec_{\beta} z = x \prec_{\alpha\beta} (y \prec_{\beta} z + y \succ_{\alpha} z + y \cdot z),$$

$$(x \succ_{\alpha} y) \prec_{\beta} z = x \succ_{\alpha} (y \prec_{\beta} z),$$

$$(x \prec_{\beta} y + x \succ_{\alpha} y + x \cdot y) \succ_{\alpha\beta} z = x \succ_{\alpha} (y \succ_{\beta} z),$$

$$(x \succ_{\alpha} y) \cdot z = x \succ_{\alpha} (y \cdot z),$$

$$(x \prec_{\alpha} y) \cdot z = x \cdot (y \succ_{\alpha} z),$$

$$(x \cdot y) \prec_{\alpha} z = x \cdot (y \prec_{\alpha} z),$$

$$(x \cdot y) \cdot z = x \cdot (y \cdot z).$$

For $n \ge 1$, let

$$T_n := T_{n,X,\Omega} := T_n(X,\Omega) \cap \{Schröder trees\}.$$

Example 10
$$T_{1} = \left\{ \begin{array}{c} x \\ \end{array} \middle| x \in X \right\},$$

$$T_{2} = \left\{ \begin{array}{c} x \\ \end{array} \middle| x \in X \right\},$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \in X \right\},$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \end{array} \middle| x \\ \left[x \in X \right],$$

$$T_{3} = \left\{ \begin{array}{c} x \\ \beta \\ \alpha \\ \end{array} \middle| x \\ \right\}$$

Denote by

$$DT(X,\Omega) := \bigoplus_{n\geq 1} \mathbf{k} T_n$$

The grafting \bigvee of $T^{(i)}, 1 \leq i \leq k$ over (x_1, \ldots, x_k) and $(\alpha_0, \ldots, \alpha_k)$ is

$$T = \bigvee_{x_1,\dots,x_k}^{k+1;\,\alpha_0,\dots,\alpha_k} (T^{(0)},\dots,T^{(k)})$$

Example 11

$$\bigvee_{u,v}^{3;\alpha,\beta,\gamma} \left(\bigvee_{v}^{x} , \bigvee_{v}^{y} , \bigvee_{v}^{z} \right) = \bigcap_{\alpha}^{\alpha} \bigcup_{v}^{\beta} \bigcap_{\gamma}^{z} .$$

Definition

Let X be a set and let Ω be a semigroup. Define binary operations

$$\prec_{\omega}, \succ_{\omega}, \cdot : \left(\mathrm{DT}(X, \Omega) \otimes \mathrm{DT}(X, \Omega) \right) \oplus \left(\mathbf{k} | \otimes \mathrm{DT}(X, \Omega) \right)$$
$$\oplus \left(\mathrm{DT}(X, \Omega) \otimes \mathbf{k} | \right) \to \mathrm{DT}(X, \Omega), \text{ for } \omega \in \Omega$$

recursively on dep(T) + dep(U) by

- 2 Let

$$T = \bigvee_{x_1,...,x_m}^{m+1; \alpha_0,...,\alpha_m} (T^{(0)},...,T^{(m)}) \in T_m, \ U = \bigvee_{y_1,...,y_n}^{n+1; \beta_0,...,\beta_n} (U^{(0)},...,U^{(n)}) \in T_n,$$

Definition

$$T \prec_{\omega} U := \bigvee_{x_{1}, \dots, x_{m-1}, x_{m}}^{m+1; \alpha_{0}, \dots, \alpha_{m-1}, \alpha_{m}\omega} (T^{(0)}, \dots, T^{(m-1)}, T^{(m)})$$

$$\succ_{\alpha_{m}} U + T^{(m)} \prec_{\omega} U + T^{(m)} \cdot U),$$

$$T \succ_{\omega} U := \bigvee_{y_{1}, y_{2}, \dots, y_{n}}^{n+1; \omega\beta_{0}, \beta_{1}, \dots, \beta_{n}} (T \succ_{\omega} U^{(0)} + T \prec_{\beta_{0}} U^{(0)})$$

$$+ T \cdot U^{(0)}, U^{(1)}, \dots, U^{(n)}),$$

$$T \cdot U := \bigvee_{x_{1}, \dots, x_{m-1}, x_{m}, y_{1}, \dots, y_{n}}^{m+n+1; \alpha_{0}, \dots, \alpha_{m-1}, \alpha_{m}\beta_{0}, \beta_{1}, \dots, \beta_{n}} (T^{(0)}, \dots, T^{(m-1)}, T^{(m)} \succ_{\alpha_{m}} U^{(0)} + T^{(m)} \prec_{\beta_{0}} U^{(0)} + T^{(m)} \cdot U^{(0)},$$

$$U^{(1)}, \dots, U^{(n)}).$$

Remark

Note that $|\prec_{\omega}|,|\succ_{\omega}|$ and $|\cdot|$ are not defined. We employ the convention that

$$|\prec_1| + |\succ_1| + |\cdot| := |,$$

and

$$|\succ_1 T := T \prec_1 | := T$$
 and $|\prec_1 T := T \succ_1 | := 0$.

Example 12

Let
$$T = X$$
, $U = X$ with $x, y, z \in X$ and $\alpha \in \Omega$. For $\beta \in \Omega$,

$$=\bigvee_{x}^{2;\,\alpha,\beta}\left(\bigvee_{y},\bigvee_{z}\right)=\bigvee_{\alpha}^{x}\bigwedge_{\beta}^{z}.$$

$$U \cdot T = = \bigvee_{z,x}^{3; 1,\alpha,1} \left(|,| \succ_1 \right) + | \prec_{\alpha} \right) + | \cdot \bigvee_{z,x}^{y} + | \cdot \bigvee_{z,x}^{y} , |$$

$$= \bigvee_{z,x}^{3; 1,\alpha,1} \left(|,| \searrow_{z}^{y} , | \right) = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y} + | \rangle_{z}^{y} + | \rangle_{z}^{y} = \bigvee_{z,x}^{y} \left(|,| \searrow_{z}^{y} , | \rangle_{z}^{y} + | \rangle_{z}^{y$$

Let $j: X \to \mathrm{DT}(X,\Omega)$ be the natural embedding map defined by $j(x) = \bigvee^x \text{ for } x \in X.$

Theorem (Z.-Gao-Manchon, Algebr. Represent. Theory)

Let X be a set and let Ω be a semigroup. Then $(\mathrm{DT}(X,\Omega),(\prec_{\omega},\succ_{\omega})_{\omega\in\Omega},\cdot)$, together with the map j, is the free tridendriform family algebra on X.

Lemma (Z.-Gao, Pacific J. Math.)

Let X be a set and let Ω be a semigroup. The Rota-Baxter family algebra $(\mathbf{k}\mathcal{T}(X,\Omega),\diamond,(B_\omega^+)_{\omega\in\Omega})$ of weight $\mathbf{1}$ induces a tridendriform family algebra $(\mathbf{k}\mathcal{T}(X,\Omega),(\prec_\omega',\succ_\omega')_{\omega\in\Omega},.')$, where

$$T \prec_{\omega}' U := T \diamond B_{\omega}^+(U), T \succ_{\omega}' U := B_{\omega}^+(T) \diamond U \text{ and } T \cdot U := T \diamond U, \text{ for } T, U \in \mathcal{T}(X, \Omega).$$

Theorem (Z.-Gao-Manchon, Algebr. Represent. Theory)

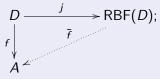
Let X be a set and let Ω be a semigroup. The free tridendriform family algebra $(\mathrm{DT}(X,\Omega),(\prec_\omega,\succ_\omega)_{\omega\in\Omega},\cdot)$ on X is a tridendriform family subalgebra of the free Rota-Baxter family algebra $(\mathbf{k}\mathcal{T}(X,\Omega),\diamondsuit,(B^+_\omega)_{\omega\in\Omega})$ of weight 1.

- 1 Free Rota-Baxter family algebras (RBFA)
 - Typed angularly decorated planar rooted trees
 - A multiplication on $kT(X, \Omega)$
 - Construction of free RBFA
- 2 Embedding free DFAs (resp.TFAs) into free RBFAs
 - Embedding free DFAs into free RBFAs
 - Embedding free TFAs into free RBFAs
 - Universal enveloping algebras of (tri)dendriform family algebras

Universal enveloping algebras

Definition

Let D be a DFA (resp. TDFA). A universal enveloping Rota-Baxter family algebra of weight λ of D satisfies the following commutative diagram:



The pair $(\mathsf{RBF}(D),j)$ is the universal enveloping RBFA of weight λ of D if

- j is a DFA (resp. TDFA) morphism (embedding map).
- A is any RBFA of weight λ, f is any DFA (resp. TDFA) morphism.
- \exists ! RBFA morphism \bar{f} .

Universal enveloping algebras

Let $j: X \to \mathrm{DD}(X,\Omega)$ be the natural embedding map defined by $j(x) = \bigvee_{x \to \infty} for \ x \in X.$

Theorem (Z.-Gao-Manchon, Algebr. Represent. Theory)

The pair $(\mathbf{k}\mathcal{T}(X,\Omega),j)$ is the universal enveloping weight 0 Rota-Baxter family algebra of the free dendriform family algebra $\mathrm{DD}(X,\Omega)$, satisfying the following commutative diagram:

$$DD(X, \Omega) \xrightarrow{j} \mathbf{k} \mathcal{T}(X, \Omega)$$

$$\downarrow g \qquad \qquad \bar{g} \qquad \qquad \bar{g}$$

$$A \stackrel{\bar{g}}{\longrightarrow} A$$

Universal enveloping algebras of (tri)dendriform family algebras

Let $j: X \to \mathrm{DT}(X,\Omega)$ be the natural embedding map defined by $j(x) = \bigvee_{x \in X} for \ x \in X.$

Theorem (Z.-Gao-Manchon, Algebr. Represent. Theory)

The pair $(\mathbf{k}T(X,\Omega),\lambda^{-1}j)$ is the universal enveloping weight λ Rota-Baxter family algebra of the free tridendriform family algebra $\mathrm{DT}(X,\Omega)$, satisfying the following commutative diagram:

$$DT(X,\Omega) \xrightarrow{\lambda^{-1}j} \mathbf{k} \mathcal{T}(X,\Omega)$$

$$\downarrow f \qquad \qquad \bar{f} \qquad \qquad \bar{f}$$

$$A \stackrel{\overline{f}}{\longrightarrow} A$$

Thank you for your attention!